

Fermat Quotients and the Polynomial Time Discrete Log Algorithm for Anomalous Elliptic Curves

by

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1. Introduction

Let p be a prime and F_p the finite field with p elements. An elliptic curve over F_p is said to be anomalous if the number of F_p rational point is exactly p . In this paper, we give an $O((\log p)^3)$ algorithm for discrete logarithm problem for an anomalous elliptic curve over a prime field. Our method may be considered as an elliptic curve version of Fermat quotient. For an integer a prime to p , the Fermat quotient $L_p(a)$ is defined to be $\frac{a^{p-1}-1}{p} \pmod{p} \in F_p$. If integers a, b are prime to p , then $L_p(ab) = L_p(a) + L_p(b)$ in F_p . Since L_p is not well defined as a function over F_p , we cannot solve discrete logarithm over F_p^\times (note also $\#F_p^\times = p-1$ is prime to $\text{char}(F_p) = p$). However, it is not so strange to expect its elliptic curve analogue is applicable to discrete log problem for an anomalous elliptic curves. We do this actually. In Sect. 2, we make a review on the Fermat quotient which illustrates our idea and historical background. We also discuss its relation to the discrete log problem on $(\mathbf{Z}/p^r\mathbf{Z})^\times$ with $p \geq 3$ and $r \geq 2$. This is not new, but does not seem to be stated explicitly. Let \tilde{E} be an anomalous curve over F_p and E its lifting to \mathbf{Z} . In Sect. 3, we construct an F_p -valued group homomorphism λ_E on the group $\tilde{E}(F_p)$ of F_p rational points of \tilde{E} for $p \geq 7$. Let $\alpha \in \tilde{E}(F_p) - \{\mathcal{O}\}$. Roughly speaking, we obtain $\lambda_E(\alpha)$ by viewing $pu(\alpha)$ in the formal group of E , where u is any lifting $\tilde{E}(F_p) \rightarrow E(\mathbf{Q}_p)$. Corollary 3.6 describes the detailed algorithm to compute λ_E . Our algorithm uses only arithmetic operations in $\mathbf{Z}/p^2\mathbf{Z}$ and F_p . This makes implementation and running time analysis simple. Then we study how to choose E so that λ_E is non-zero. Summing up, we can solve the discrete log problem of \tilde{E} . We discuss the cryptographic implication of our result in Sect. 4.

After the works on the discrete log problem for anomalous elliptic curves over prime fields were completed, the authors were informed that Dr. N. Smart has independently obtained same results [21] at the similar time.

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theory and its related topics" held at RIMS, Kyoto University on Oct. 27–31, 1997. Later (precisely, on Nov. 3, 1997), the authors learned that Semaev [17] obtained (in 1995) an polynomial time discrete log algorithm of p -torsion points of an elliptic curves over F_p and that Rück [16] generalized it for curves of arbitrary genus. However, the method of Semaev and Rück which is algebraic geometric is quite different to ours and that of Smart. For the comparison of these two algorithms, see Voloch [22].

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Notation

Rings are always assumed to be commutative and unitary. Let R be a ring. We denote the unit group (the set of invertible elements) by R^\times . For a prime p , the finite field with p elements is denoted by F_p . For $a \in \mathcal{Q}$, we put

$$\text{ord}_p a := \begin{cases} r & \left(a = p^r \frac{v}{u}, u, v \in \mathbf{Z} - p\mathbf{Z} \right), \\ \infty & (a = 0). \end{cases}$$

We denote the p -adic number field and the ring of p -adic integers by \mathcal{Q}_p and \mathbf{Z}_p , respectively. By definition, \mathcal{Q}_p is the completion of \mathcal{Q} with respect to the metric induced from ord_p . The function ord_p uniquely extends to a continuous function $\mathcal{Q}_p \rightarrow \mathbf{Z} \cup \{\infty\}$. The following formulas are well known.

$$\begin{aligned} \text{ord}_p(xy) &= \text{ord}_p x + \text{ord}_p y & (x, y \in \mathcal{Q}_p) \\ \text{ord}_p(x+y) &\geq \min(\text{ord}_p x, \text{ord}_p y) \\ \text{ord}_p(x+y) &= \min(\text{ord}_p x, \text{ord}_p y) & \text{for } \text{ord}_p x \neq \text{ord}_p y \\ \mathbf{Z}_p &= \{x \in \mathcal{Q}_p : \text{ord}_p x \geq 0\} \\ \mathbf{Z}_p^\times &= \{x \in \mathcal{Q}_p : \text{ord}_p x = 0\} \end{aligned}$$

For introductory explanation of p -adic numbers, see e.g. Cassels [2], Serre [18].

2. Fermat quotient

Let p be a prime. In 1828, Abel [1] proposed the following problem: Can a number $a^{p-1} - 1$ be divisible by p^2 , where p is a prime and a is an integer, $1 \leq a < p$? Note we always have $n^{p-1} - 1 \equiv 0 \pmod{p}$ for $n \in \mathbf{Z} - p\mathbf{Z}$ due to the Fermat's little theorem. Dickson [3, p. 105] states that G. Eisenstein noted the following formulas in 1850:

$$\begin{cases} L_p(ab) = L_p(a) + L_p(b) \\ L_p(a+pc) = L_p(a) - ca^{-1} \end{cases} \quad (2.1)$$

for $a, b \in \mathbf{Z} - p\mathbf{Z}$, $c \in \mathbf{Z}$ where

$$L_p(a) := \frac{a^{p-1} - 1}{p} \bmod p \in \mathbf{F}_p \quad (2.2)$$

and a^{-1} is the inverse of a in \mathbf{F}_p . We call L_p the *Fermat quotient*. Lerch [9, (27)] noticed generalization of the Fermat quotient to a composite modulus $m \in \mathbf{N}$. For $a, m \in \mathbf{Z}$ with $\gcd(a, m) = 1$, put $L_m(a) := \frac{a^{\varphi(m)} - 1}{m} \bmod m \in \mathbf{Z}/m\mathbf{Z}$. Then,

$$\begin{cases} L_m(ab) = L_m(a) + L_m(b) \\ L_m(a + mc) = L_m(a) + \varphi(m)ca^{-1} \end{cases} \quad (2.3)$$

where $\gcd(a, m) = \gcd(b, m) = 1$ and $c \in \mathbf{Z}$. For more details, see Dickson [3, Chap. 4]. Concerning Abel's original problem, Jacobi [7] observed $a^{p-1} \equiv 1 \bmod p^2$ has only four solutions $(a, p) = (3, 11), (9, 11), (14, 29), (18, 37)$ for $p \leq 37$. As far as the authors know, it is open whether there are infinitely many primes satisfying $a^{p-1} \equiv 1 \bmod p^2$ for a fixed a . Ihara [5] considers this problem from modern number theoretic point of view. Ribenboim [15, Chap. 5.III] states that Crandall, Dilcher and Pomerance verified only $p = 1093$ and 3511 satisfy $2^{p-1} \equiv 1 \bmod p^2$ for $p < 4 \times 10^{12}$.

The phenomena concerning (2.2) may be paraphrased as follows: Let G be a finite group of order n consisting of $\bmod p$ objects. For $g \in G$, consider g^n in $\bmod p^2$ (via some "lifting"), which should be p -adically "close" to the identity of G . The difference between them may involve interesting information on n . We do this for an anomalous elliptic curve defined over \mathbf{F}_p in the next section.

In order to observe relation between Fermat quotient and discrete logarithm, let us consider the discrete log problem $\bmod p^r$ for $p \geq 3$ and $r \geq 2$. Let $\omega \in \mathbf{Z}$ be a primitive element of $\mathbf{Z}/p^2\mathbf{Z}$. Then, as is well known, ω is a primitive element of $\mathbf{Z}/p^r\mathbf{Z}$ for all $r \geq 1$. For $\alpha \in (\mathbf{Z}/p^r\mathbf{Z})^\times$, we want $n \in \mathbf{Z}/p^{r-1}(p-1)\mathbf{Z}$ satisfying

$$\alpha \equiv \omega^n \bmod p^r. \quad (2.4)$$

For $a \in \mathbf{Z} - p\mathbf{Z}$, note $L_{p^r}(a + p^r) = L_{p^r}(a) - p^{r-1}a^{-1}$ by (2.3). Hence L_{p^r} induces a well defined map $(\mathbf{Z}/p^r\mathbf{Z})^\times \rightarrow \mathbf{Z}/p^{r-1}\mathbf{Z}$, which is again denoted by L_{p^r} . By (2.3), $L_{p^r}(\alpha) \equiv nL_{p^r}(\omega) \bmod p^{r-1}$. On the other hand, $\omega^{p-1} \equiv 1 + pL_p(\omega) \bmod p^2$. Then $\omega^{p-1} \not\equiv 1 \bmod p^2$ implies $L_p(\omega) \in \mathbf{F}_p^\times$, while $\omega^{(p-1)p^{r-1}} \equiv (1 + pL_p(\omega))^{p^{r-1}} \bmod p^{r+1}$. Since $p \geq 3$, we see $\omega^{(p-1)p^{r-1}} \equiv 1 + p^r L_p(\omega) \bmod p^{r+1}$ (cf. Ireland and Rosen [6, Chap. 4, Sect. 1, Lemma 3 and Corollary 1]). Therefore, $L_{p^r}(\omega) \equiv L_p(\omega) \bmod p$, i.e., $L_{p^r}(\omega) \in (\mathbf{Z}/p^r\mathbf{Z})^\times$. Consequently, we obtain

$$n \equiv \frac{L_{p^r}(\alpha)}{L_{p^r}(\omega)} \bmod p^{r-1}. \quad (2.5)$$

On the other hand, $\bmod p$ of (2.4) yields $\alpha \equiv \omega^n \bmod p$. Any discrete log algorithm for \mathbf{F}_p gives $k \in \mathbf{Z}/(p-1)\mathbf{Z}$ such that

$$n \equiv k \pmod{p-1}. \quad (2.6)$$

Using Chinese remainder theorem, we obtain n from (2.5) and (2.6). Let $T(p)$ be a time complexity for discrete logarithm in \mathbf{F}_p . The time complexity of above method is $O(T(p) + (\log p^r)^2 \log r)$, which runs faster than Pohlig-Hellman algorithm [14]. We note that we only need $a^{p^r - p^{r-1}} \pmod{p^{2r-1}}$ to compute $L_{p^r}(a) \pmod{p^{r-1}}$.

3. Discrete log problems on anomalous elliptic curves

Let p be a prime. Let

$$\tilde{E}: y^2 + \tilde{a}_1 xy + \tilde{a}_3 y = x^3 + \tilde{a}_2 x^2 + \tilde{a}_4 x + \tilde{a}_6$$

be an elliptic curve defined over \mathbf{F}_p . We say \tilde{E} is *anomalous* if $\#\tilde{E}(\mathbf{F}_p)$ (including the point at infinity) is exactly p . Mazur [10] studied such curves in extensively sophisticated setting. Let E be an lifting of \tilde{E} to \mathbf{Z} . In other words, choose $a_i \in \mathbf{Z}$ satisfying $a_i \pmod{p} = \tilde{a}_i$ for $i = 1 \sim 4, 6$ and define E by

$$E: y^2 + a_1 xy + a_3 y = x^3 + a_2 x^2 + a_4 x + a_6.$$

We denote the reduction map $E(\mathbf{Q}_p) \rightarrow \tilde{E}(\mathbf{F}_p)$ by π . The identity element of an elliptic curve with respect to its group structure is denoted by \mathcal{O} . For any \mathbf{Z} -algebra R , we put

$$E(R) := \{(x : y : 1) \in \mathbf{P}^2(R) : y^2 + a_1 xy + a_3 y = x^3 + a_2 x^2 + a_4 x + a_6\} \cup \{(0 : 1 : 0)\}.$$

We note $E(R)$ is *not* necessarily a group if R is not a field. We identify $(X : Y : 1) \in \mathbf{P}^2(R)$ with $(X, Y) \in A^2(R)$. By abuse of notation, we write $(0 : 1 : 0)$ as \mathcal{O} . Let \mathcal{E} be the Formal group associated to E . (See Silverman [20, Chap. 4] for its definition and its basic properties.) We have a following isomorphism \mathcal{L} :

$$\mathcal{L}: \text{Ker } \pi \xrightarrow{\psi} \mathcal{E}(p\mathbf{Z}_p) \xrightarrow{\log_{\mathcal{E}}} p\mathbf{Z}_p \quad (3.1)$$

where $\psi(x : y : z) = x/y$ and

$$\begin{aligned} \log_{\mathcal{E}}(t) := & t - \frac{a_1}{2} t^2 + \frac{a_1^2 + a_2}{3} t^3 - \frac{a_1^3 + 2a_1 a_2 + a_3}{4} t^4 \\ & + \frac{a_1^4 + 3a_1^2 a_2 + 6a_1 a_3 + a_2^2 + 2a_4}{5} t^5 - \dots \end{aligned} \quad (3.2)$$

is a formal logarithm of \mathcal{E} . In what follows, we always assume \tilde{E} is anomalous.

LEMMA 3.1. $pE(\mathbf{Q}_p) \subset \text{Ker } \pi$.

Proof. Let $A \in E(\mathbf{Q}_p)$. Since π is a group homomorphism by Silverman [20, Chap. 7, proof of Prop. 2.1], anomaly of \tilde{E} implies $\pi(pA) = p\pi(A) = \mathcal{O}$. \square

THEOREM 3.2. Let u be any lifting from $\tilde{E}(\mathbf{F}_p)$ to $E(\mathbf{Q}_p)$, i.e. $\pi \circ u = \text{id}_{\tilde{E}(\mathbf{F}_p)}$. Let λ_E be a composition of the following maps

$$\lambda_E: \tilde{E}(\mathbf{F}_p) \xrightarrow{u} E(\mathbf{Q}_p) \xrightarrow{h_p} \text{Ker } \pi \xrightarrow{\mathcal{L}} p\mathbf{Z}_p \xrightarrow{\text{mod } p^2} p\mathbf{Z}_p/p^2\mathbf{Z}_p \cong \mathbf{F}_p. \quad (3.3)$$

Here h_p is multiplication by p . Then λ_E is a group homomorphism which is independent of choice of u . Moreover, λ_E is either a zero map or an isomorphism.

Proof. Let $\alpha, \beta \in \tilde{E}(\mathbf{F}_p)$ and put $\Delta := u(\alpha) + u(\beta) - u(\alpha + \beta)$. Since π is a group homomorphism and u is a lifting, we see $\pi(\Delta) = \mathcal{O}$, i.e. $\Delta \in \text{Ker } \pi$. By (3.1) there is $t_0 \in \mathbf{Z}_p$ satisfying $\mathcal{L}(\Delta) = pt_0$. Therefore $\mathcal{L}(h_p(\Delta)) = p^2t_0 \in p^2\mathbf{Z}_p$. We consequently have $F(\Delta) = 0$, where $F := (\text{mod } p^2) \circ \mathcal{L} \circ h_p$. Since F is a homomorphism, $F(u(\alpha)) + F(u(\beta)) = F(u(\alpha + \beta))$, i.e., λ_E is a homomorphism. Let v be another lifting $\tilde{E}(\mathbf{F}_p) \rightarrow E(\mathbf{Q}_p)$. Then $\pi(u(\alpha) - v(\alpha)) = \pi(u(\alpha)) - \pi(v(\alpha)) = \mathcal{O}$ for any $\alpha \in \tilde{E}(\mathbf{F}_p)$. Hence we have $u(\alpha) - v(\alpha) \in \text{Ker } \pi$. By the same argument as above we see that $F(u(\alpha)) = F(v(\alpha))$ and that λ_E is independent of choice of u . To show the last statement, simply note $\tilde{E}(\mathbf{F}_p)$ has no nontrivial proper subgroup. Therefore either $\text{Ker } \lambda_E$ is $\tilde{E}(\mathbf{F}_p)$, which implies λ_E is a zero map, or $\text{Ker } \lambda_E$ is trivial, which implies λ_E is injective. In the latter case, λ_E is also surjective since ${}^*\tilde{E}(\mathbf{F}_p) = p = {}^*\mathbf{F}_p$. Hence λ_E is an isomorphism. \square

REMARK 3.3. Although λ_E is independent of choice of u , it depends on the choice of E . Cf. Theorem 3.7.

COROLLARY 3.4. Let \tilde{E} be an anomalous elliptic curve defined over \mathbf{F}_p and E its lifting to \mathbf{Z} . Then, the following conditions on E are equivalent:

- (i) The map λ_E is a zero map.
- (ii) There exists $\alpha \in \tilde{E}(\mathbf{F}_p) - \{\mathcal{O}\}$ satisfying $\lambda_E(\alpha) = 0$.
- (iii) There exists a p -torsion point belonging to $E(\mathbf{Z}_p) - \{\mathcal{O}\}$.

Proof. The implication (i) \rightarrow (ii) is trivial. (ii) \rightarrow (i): Let $\lambda_E(\alpha) = 0$ with $\alpha \in \tilde{E}(\mathbf{F}_p) - \{\mathcal{O}\}$ and let $\beta \in \tilde{E}(\mathbf{F}_p)$ be arbitrary. Since $\tilde{E}(\mathbf{F}_p)$ is a cyclic group of order p , α is its generator. Hence there exists $n \in \mathbf{N}$ satisfying $\beta = n\alpha$. Therefore $\lambda_E(\beta) = n\lambda_E(\alpha) = 0$.

(ii) \rightarrow (iii): Let $\alpha \in \tilde{E}(\mathbf{F}_p) - \{\mathcal{O}\}$ with $\lambda_E(\alpha) = 0$. Then there is $t_1 \in \mathbf{Z}_p$ satisfying $\mathcal{L}(pu(\alpha)) = p^2t_1$. Put $B := \mathcal{L}^{-1}(pt_1) \in \text{Ker } \pi$. Since \mathcal{L} is an isomorphism, $pB = pu(\alpha)$. Letting $A := u(\alpha) - B$, we see A is a p -torsion point belonging to $E(\mathbf{Z}_p) - \{\mathcal{O}\}$.

(iii) \rightarrow (ii): Assume $A \in E(\mathbf{Z}_p) - \{\mathcal{O}\}$ and $pA = \mathcal{O}$. Put $\alpha := \pi(A)$. Then

$$\lambda_E(\alpha) = ((\text{mod } p^2) \circ \mathcal{L})(pA) = \mathcal{O}.$$

On the other hand, $\alpha \neq \mathcal{O}$ by $A \in E(\mathbf{Z}_p) - \{\mathcal{O}\}$. \square

In what follows let $p \geq 5$ for simplicity. Then without loss of generality, we may assume $\tilde{a}_1 = \tilde{a}_2 = \tilde{a}_3 = 0$ (in \mathbf{F}_p) and $a_1 = a_2 = a_3 = 0$ (in \mathbf{Z}).

THEOREM 3.5. Let $p \geq 5$ be a prime. Let $\alpha \in \tilde{E}(\mathbf{F}_p) - \{\mathcal{O}\}$. Assume $A := (x_1, y_1) \in E(\mathbf{Z}_p)$ satisfies $\pi(A) = \alpha$. For those $n \in \mathbf{N}$ such that $nA \neq \mathcal{O}$, put $(x_n, y_n) := nA$. If λ_E is a non-zero map, we have the following:

- (i) $nA \in E(\mathbf{Z}_p) - \{\mathcal{O}\}$ for $1 \leq n < p$,

- (ii) $x_n \not\equiv x_m \pmod p$ for $1 \leq n < m < p$ with $n+m \neq p$,
 (iii) $y_{p-1} - y_1 \in \mathbf{Z}_p^\times$, $\frac{x_{p-1} - x_1}{p} \in \mathbf{Z}_p^\times$ and $\lambda_E(\alpha) = \frac{x_{p-1} - x_1}{p(y_{p-1} - y_1)} \pmod p$.

Proof. (i) First of all, we note $nA \neq \mathcal{O}$ for $1 \leq n < p$. Indeed, $nA = \mathcal{O}$ implies $n\alpha = \pi(nA) = \mathcal{O}$. Since \tilde{E} is anomalous, we obtain $\alpha = \mathcal{O}$, which is a contradiction. So we have only to prove $nA \in E(\mathbf{Z}_p)$. For $n=1$, this is a part of assumptions. For $n=2$, we note

$$y_1 \not\equiv 0 \pmod p. \quad (3.4)$$

Otherwise, we obtain $2\alpha = 2\pi(A) = 2(x_1 \pmod p, 0) = \mathcal{O}$, which contradicts to $\alpha \neq \mathcal{O}$ since \tilde{E} is anomalous. Therefore, we have $y_1 \in \mathbf{Z}_p^\times$. Addition formula of E yields

$$x_2 = c_2^2 - 2x_1, \quad y_2 = -c_2x_2 - d_2,$$

where

$$c_2 = \frac{3x_1^2 + a_4}{2y_1}, \quad d_2 = \frac{-x_1^3 + a_4x_1 + 2a_6}{2y_1}.$$

Since $y_1 \in \mathbf{Z}_p^\times$, we see $x_2, y_2 \in \mathbf{Z}_p$ and (i) holds for $n=2$. For $3 \leq n < p$, we use induction on n . Suppose $A, (n-1)A \in E(\mathbf{Z}_p) - \{\mathcal{O}\}$. Especially,

$$\pi(A) = (x_1 \pmod p, y_1 \pmod p)$$

$$\pi((n-1)A) = (x_{n-1} \pmod p, y_{n-1} \pmod p).$$

Assuming $x_1 \equiv x_{n-1} \pmod p$, we obtain $\pi(A) = \pm \pi((n-1)A)$, i.e., $n\alpha = \mathcal{O}$ or $(n-2)\alpha = \mathcal{O}$. By the anomaly of \tilde{E} , we have $\alpha = \mathcal{O}$, which is again a contradiction. Therefore we obtain $x_1 \not\equiv x_{n-1} \pmod p$ and, a fortiori, $x_1 \neq x_{n-1}$. Then,

$$x_n = c_n^2 - x_1 - x_{n-1}, \quad y_n = -c_n^3 + c_n(x_1 + x_{n-1}) - d_n, \quad (3.5)$$

where

$$c_n = \frac{y_{n-1} - y_1}{x_{n-1} - x_1}, \quad d_n = \frac{y_1x_{n-1} - y_{n-1}x_1}{x_{n-1} - x_1}. \quad (3.6)$$

By $x_{n-1} \neq x_1 \pmod p$, we see $x_{n-1} - x_1 \in \mathbf{Z}_p^\times$ and hence $c_n, d_n \in \mathbf{Z}_p$. Therefore $x_n, y_n \in \mathbf{Z}_p$.

The proof of (ii) is similar. Assume $x_n \equiv x_m \pmod p$. Then $\pi(nA) = \pm \pi(mA)$, i.e., $(m \pm n)\alpha = \mathcal{O}$, which yields a contradiction as in (i).

(iii): Since $\lambda_E \neq 0$, we see $pA \neq \mathcal{O}$ (cf. Corollary 3.4). Note (3.5) and (3.6) hold for $n=p$. Let $(x_p, y_p) := pA$. Then $\pi((x_p : y_p : 1)) = \mathcal{O} = (0 : 1 : 0)$ implies $\text{ord}_p y_p < 0$ and $\text{ord}_p x_p > \text{ord}_p y_p$. By (ii), we have $A, (p-1)A \in E(\mathbf{Z}_p)$. Let $s := \text{ord}_p c_p$. Assume $s \geq 0$, i.e. $c_p \in \mathbf{Z}_p$. Using

$$d_p = y_1 - x_1 c_p, \quad (3.7)$$

we have $d_p \in \mathbf{Z}_p$, and hence $y_p \in \mathbf{Z}_p$, a contradiction. So, s must be negative. Then,

by (3.5), we see

$$\text{ord}_p x_p = 2s . \quad (3.8)$$

By (3.7), we also obtain

$$\begin{aligned} \text{ord}_p d_p &\geq \min(\text{ord}_p y_1, \text{ord}_p x_1 + \text{ord}_p c_p) \\ &\geq \text{ord}_p c_p = s . \end{aligned}$$

Moreover $\text{ord}_p(c_p(x_1 + x_{p-1})) \geq s$, while $\text{ord}_p c_p^3 = 3s < s$. Hence

$$\text{ord}_p y_p = 3s . \quad (3.9)$$

Therefore, $\text{ord}_p \psi(pA) = \text{ord}_p(x_p/y_p) = -s > 0$. Using (3.2), we see $\text{ord}_p \mathcal{L}(pA) = -s$.

By assumption $\lambda_E(A) \neq 0$. So $\text{ord}_p \mathcal{L}(pA) = 1$. Summing up, we obtain $s = -1$, $\frac{x_p}{py_p} \in$

\mathbf{Z}_p^\times and $\lambda_E(A) = \frac{x_p}{py_p} \bmod p$. By the anomaly of \tilde{E} , we see $\pi((p-1)A) = -\pi(A)$ and hence $y_{p-1} \equiv -y_1 \bmod p$. Therefore, $y_{p-1} - y_1 \equiv -2y_1 \not\equiv 0 \bmod p$. So, we have proved $y_{p-1} - y_1 \in \mathbf{Z}_p^\times$. Since $\text{ord}_p c_p = -1$, we obtain

$$\frac{x_{p-1} - x_1}{p} \in \mathbf{Z}_p^\times . \quad (3.10)$$

Let $\hat{x} := p^2 x_p$ and $\hat{y} := p^3 y_p$. By $s = -1$, (3.8) and (3.9), we see $\hat{x}, \hat{y} \in \mathbf{Z}_p^\times$. Hence $\lambda_E(A) = \frac{\hat{x}}{\hat{y}} \bmod p = \frac{\hat{x} \bmod p}{\hat{y} \bmod p}$. Note $pc_p \in \mathbf{Z}_p^\times$ since $s = -1$. Therefore,

$$\begin{aligned} \hat{x} \bmod p &= (p^2 c_p^2 - p^2(x_1 + x_{p-1})) \bmod p \\ &= (pc_p)^2 \bmod p \end{aligned}$$

and

$$\begin{aligned} \hat{y} \bmod p &= -p^3 c_p^3 + (pc_p)p^2(x_1 + x_{p-1}) - p^3 d_p \bmod p \\ &= -(pc_p)^3 \bmod p . \end{aligned}$$

Consequently,

$$\lambda_E(\alpha) = \frac{(pc_p)^2 \bmod p}{-(pc_p)^3 \bmod p} = \left(-\frac{1}{p} \frac{x_{p-1} - x_1}{y_{p-1} - y_1} \right) \bmod p . \quad (3.11)$$

This completes the proof. \square

COROLLARY 3.6. *Let $p \geq 5$. Let*

$$\tilde{E}: y^2 = x^3 + \tilde{a}_4 x + \tilde{a}_6$$

be an anomalous elliptic curve defined over \mathbf{F}_p . Choose integers a_4, a_6 satisfying $a_4 \bmod p = \tilde{a}_4$ and $a_6 \bmod p = \tilde{a}_6$. Define an elliptic curve E by

$$E: y^2 = x^3 + a_4 x + a_6 .$$

Let λ_E be a homomorphism defined by (3.3). Then, the following procedure computes $\lambda_E(\alpha)$ for $\alpha := (s, t) \in \tilde{E}(\mathbf{F}_p) - \{\mathcal{O}\}$ with $O((\log p)^3)$ time complexity.

- (i) Find $A := (X_1, Y_1) \in E(\mathbf{Z}/p^2\mathbf{Z})$ satisfying $X_1 \bmod p = s$ and $Y_1 \bmod p = t$.
- (ii) Compute $(X_{p-1}, Y_{p-1}) := (p-1)A \in E(\mathbf{Z}/p^2\mathbf{Z})$ by using elliptic curve addition.
- (iii) If $X_{p-1} \neq X_1$, then

$$\lambda_E(\alpha) = \left(\frac{X_{p-1} - X_1}{p} \bmod p \right) ((Y_{p-1} - Y_1) \bmod p)^{-1}.$$

Otherwise $\lambda_E(\alpha) = 0$.

Proof. Note, under the same notation as in Theorem 3.5, we have only to compute $y_{p-1} - y_1 \bmod p$ and $\frac{1}{p}(x_{p-1} - x_1) \bmod p$. For (i), simply take any $X_1, y \in \mathbf{Z}/p^2\mathbf{Z}$ satisfying $X_1 \bmod p = s$ and $y \bmod p = t$. Then solve the following equation on w :

$$(y + pw)^2 = X_1^3 + a_4X_1 + a_6 \bmod p^2,$$

$$\text{i.e. } 2tw = \frac{X_1^3 + a_4X_1 + a_6 - y^2}{p} \bmod p.$$

Note the right hand side is well defined. By (3.4), we obtain $w \in \mathbf{F}_p$. Then put $Y_1 := y + pw$. See Serre [18, Chap. 2, §2.2]. Then Theorem 3.5 (ii) guarantees computation of $(p-1)A \bmod p^2$ involves only operations over $\mathbf{Z}/p^2\mathbf{Z}$. By (3.10), we see $X_{p-1} \neq X_1$ under the condition $\lambda_E \neq 0$. In this case Theorem 3.5 (iii) ensures validity of Step (iii). Otherwise λ_E must be a zero map. So, $\lambda_E(\alpha) = 0$.

The number of arithmetic operations over $\mathbf{Z}/p^2\mathbf{Z}$ involved in steps (i) and (iii) are indifferent to p or \tilde{E} . Step (ii) requires at most $2 \log_2 p$ elliptic curve addition. Summing up, $O((\log p)^3)$ is enough to compute $\lambda_E(\alpha)$. \square

THEOREM 3.7. *Let E and \tilde{E} be as in Corollary 3.6, $A := (x_1, y_1) \in E(\mathbf{Z}_p) - \{\mathcal{O}\}$. Assume*

$$3x_1^2 \not\equiv -a_4 \bmod p. \quad (3.12)$$

Define an elliptic curve E' by

$$E' : y^2 = x^3 + a'_4x + a'_6,$$

where $a'_4 := a_4 + p$ and $a'_6 := a_6 - px_1$. (Note $A \in E'(\mathbf{Z}_p)$.) Let $(x_{p-1}, y_{p-1}) := (p-1)A$ in E and $(x'_{p-1}, y'_{p-1}) := (p-1)A$ in E' . Then, we have the following:

- (i) *Either $x_{p-1} \not\equiv x'_{p-1} \bmod p^2$ or $y_{p-1} \not\equiv y'_{p-1} \bmod p^2$.*
- (ii) *Either λ_E or $\lambda_{E'}$ is non-zero.*

Proof. (i): Assume $x_{p-1} \equiv x'_{p-1} \bmod p^2$ and $y_{p-1} \equiv y'_{p-1} \bmod p^2$. For $1 \leq n \leq p-1$, let $(x_n, y_n) := nA$ in E and $(x'_n, y'_n) := nA$ in E' . Writing elliptic curve addition

formula for $(n-1)A = nA + (-A)$ in E explicitly, we have

$$x_{n-1} = \hat{c}_n^2 - x_1 - x_n, \quad y_{n-1} = -\hat{c}_n^3 + \hat{c}_n(x_1 + x_n) - \hat{d}_n$$

where

$$\hat{c}_n = \frac{y_n + y_1}{x_n - x_1}, \quad \hat{d}_n = -\frac{y_1 x_n + y_n x_1}{x_n - x_1}$$

for $n \geq 3$. The similar formulas hold for (x'_{n-1}, y'_{n-1}) . But these formulas do not contain coefficients of the Weierstrass equation. By Theorem 3.5 (ii), $x_n \equiv x'_n \pmod{p^2}$ and $y_n \equiv y'_n \pmod{p^2}$ implies $x_{n-1} \equiv x'_{n-1} \pmod{p^2}$ and $y_{n-1} \equiv y'_{n-1} \pmod{p^2}$. On the other hand, the duplication formula on E and E' yield

$$\begin{aligned} x'_2 - x_2 &= \left(\frac{3x_1^2 + a'_4}{2y_1} \right)^2 - 2x_1 - \left(\left(\frac{3x_1^2 + a_4}{2y_1} \right)^2 - 2x_1 \right) \\ &= p \frac{6x_1^2 + 2a_4 + p}{4y_1^2} \\ &\not\equiv 0 \pmod{p^2} \end{aligned}$$

by (3.12), which is a contradiction.

(ii): Assume both λ_E and $\lambda_{E'}$ are zero maps. By (3.11), $\text{ord}_p(x_{p-1} - x_1) \geq 2$ and $\text{ord}_p(x'_{p-1} - x_1) \geq 2$. Hence $x_{p-1} \equiv x'_{p-1} \pmod{p^2}$. Then,

$$\begin{aligned} y'_{p-1} - y_{p-1} &= x'_{p-1}{}^3 - x_{p-1}{}^3 + a'_4 x'_{p-1} - a_4 x_{p-1} + a'_6 - a_6 \\ &\equiv x_1(a'_4 - a_4) + (a'_6 - a_6) \pmod{p^2} \\ &\equiv 0 \pmod{p^2} \end{aligned}$$

which contradicts to (i). \square

COROLLARY 3.8. *Let $p \geq 7$. Let \tilde{E} be an anomalous elliptic curve defined over \mathbf{F}_p . Then we can solve discrete logarithm for $\tilde{E}(\mathbf{F}_p)$ with $O((\log p)^3)$ time complexity.*

Proof. Let $\alpha, \beta \in \tilde{E}(\mathbf{F}_p) - \{\mathcal{O}\}$. Since \tilde{E} is anomalous, there exists $n \in \mathbf{F}_p$ satisfying $\beta = n\alpha$. Define E as in Corollary 3.6. Since λ_E is a homomorphism by Theorem 3.2, we have $\lambda_E(\beta) = n\lambda_E(\alpha)$. In case of $\lambda_E(\alpha) \neq 0$, we obtain $n = \frac{\lambda_E(\beta)}{\lambda_E(\alpha)}$. Otherwise, λ_E is a zero map by Corollary 3.4. By Theorem 3.5 (ii), at least one of $u(\alpha)$, $u(2\alpha)$, or $u(3\alpha)$ has an x -coordinate satisfying (3.12). (Note $p \geq 7$.) Therefore we can form a non-zero $\lambda_{E'}$ by Theorem 3.7 in $O((\log p)^2)$ time complexity. Thus, $n = \frac{\lambda_{E'}(\beta)}{\lambda_{E'}(\alpha)}$. The total running time of above procedure is clearly $O((\log p)^3)$. \square

4. Concluding remarks[†]

4.1. Non-prime field case

Let p be a prime, $q := p^r$ with $r \geq 1$, and K be an unramified extension of \mathcal{Q}_p whose residue field is F_p . See Serre [19, Chap. I, Sect. 6] for explicit construction of K . Let \tilde{E} be an elliptic curve defined over F_q . Let $A \in \tilde{E}(F_q)$ and N its order, i.e., the cardinality of the cyclic group $\langle A \rangle$ generated by A . We write $N = p^e m$ where m is prime to p . We consider the discrete log problem for $\langle A \rangle$ in case of $e \geq 1$. Let $B = nA$ with $n \in \mathbb{Z}/N\mathbb{Z}$. Even if $m > 1$, we obtain $n \bmod p^e$ in polynomial time as follows. There exist integers n_0, \dots, n_{e-1} satisfying $n \equiv \sum_{i=0}^{e-1} n_i p^i \pmod{p^e}$ and $0 \leq n_i < p$. We put $A_0 := mp^{e-1}A$ and $B_0 := mp^{e-1}B$. Then $pA_0 = \mathcal{O}$ and $B_0 = n_0 A_0$. In case of a small p , we obtain n_0 by checking whether $B_0 = n_0 A_0$ for $n_0 = 0, 1, \dots, p-1$. Otherwise (at least $p \geq 7$), we use the elliptic Fermat quotient. Recall that the essential points of Sect. 3 are the following:

- (i) The order of the base point A is p .
- (ii) The formal logarithm $\log_{\mathfrak{g}}$ is defined over $\psi(\text{Ker } \pi)$.

Since K/\mathcal{Q}_p is unramified, (ii) also holds for an elliptic curve over K . Choosing a lifting E of \tilde{E} to the valuation ring of K , we obtain n_0 by $\lambda_E(B_0)/\lambda_E(A_0)$. In order to obtain n_i for $i \geq 1$, we use Pohlig-Hellman algorithm [14]. Assume we have obtained n_0, \dots, n_{k-1} . Then $B_k := mp^{e-k-1}(B - (\sum_{i=0}^{k-1} n_i p^i)A)$ satisfies $B_k = n_k A_0$, which yields n_k by the same method as above. Repeating this process, we obtain $n \bmod p^e$.

4.2. Density of vulnerable curves

It is natural to ask how often an elliptic curve is anomalous. By McKee [11, Theorem 2], the density of anomalous curves over F_p is at most $O\left(\frac{1}{\sqrt{p}} \log p \log \log p\right)$. In case of a large p , this is fairly small. For $q = p^r$, the probability of $p \mid \# \tilde{E}(F_q)$ converges to $1/(p-1)$ as $r \rightarrow \infty$ by Howe [4, Theorem 1.1].

4.3. Cryptographic implication

Since Miller [13] and Koblitz [8] independently proposed the elliptic curve cryptosystem, many works are done about this cryptosystem and a number of cryptographic apparatuses using elliptic curve discrete log problem are now commercially available. Among them, Menezes-Okamoto-Vanstone [12] showed that the discrete logarithm problem of supersingular curves is reduced to that of multiplicative group of some extensions of the base field. Since the discrete logarithm of finite field is still considered to be difficult (sub-exponential time), the elliptic curve cryptosystem is believed to be no less secure than cryptosystem based on the discrete log problem of a finite field. However, our result indicates that we should choose the elliptic curve and its base point carefully. At least we should

[†] After the release of the first version of this paper, the authors received several comments and questions on it. Some of them fit better in the previous section. However, to be fair, we describe them here.

choose a base point whose order is divisible by a large prime other than the characteristic of a base field. This automatically excludes anomalous curves. Of course, this condition is one of necessary conditions for unvulnerability, *not* a sufficient condition.

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